

## CMPE 350 - Spring 2018

### PS 6 - 19.03.18

**2.26** Show that if  $G$  is a CFG in Chomsky Normal Form, then for any string  $w \in L(G)$  of length  $n \geq 1$ , exactly  $2n - 1$  steps are required for any derivation of  $w$ .

**2.5** Give informal descriptions and state diagrams of pushdown automata for the languages in 2.4.

**2.18 a)** Let  $C$  be a context-free language and  $R$  be a regular language. Prove that the language  $C \cap R$  is context-free.

**b)** Use part a) to show that the language  $A = \{w \mid w \in \{a, b, c\}^* \text{ and contains equal number of } a\text{'s, } b\text{'s and } c\text{'s}\}$  is not a CFL.

**2.44** If  $A$  and  $B$  are languages, define  $A \diamond B = \{xy \mid x \in A \text{ and } y \in B \text{ and } |x| = |y|\}$ . Show that if  $A$  and  $B$  are regular languages, then  $A \diamond B$  is CFL.

- Prove that there are infinitely many context-free languages which are non-regular.